

Information Percolation in Large Markets

from work by

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Information Percolation

▶ Markets

- Hayek (1945)
- Wolinsky (1990)
- Golosov, Lorenzoni, Tsyvinski (2008)

▶ Social learning

- Banerjee and Fudenberg (1995)
- Acemoglu (2008)

Setting

From Duffie and Manso (AER 2007):

- ▶ A continuum of agents matched pairwise independently to other agents at mean rate r .
- ▶ Payoff relevant states: $X = \begin{cases} H & \text{with probability } \nu \\ L & \text{with probability } 1 - \nu \end{cases}$
- ▶ Agent k is endowed with $S_k = \{s_1, \dots, s_{k_n}\}$, $\{0, 1\}$ -signals that are X -conditionally independent, with

$$P(s_i = 1 | X = H) \geq P(s_i = 1 | X = L).$$

- ▶ For almost every pair j and k of agents, S_j and S_k are disjoint.
- ▶ If j and k are matched, they share endowed and previously gathered signals.

Information is Additive in Types

- ▶ For any conditional probability $p \in (0, 1)$ of the event $\{X = H\}$, we define the associated information type

$$\Theta(p) = \log \frac{(1-p)\nu}{(1-\nu)p}$$

- ▶ Result: Sharing information is additive in types. That is, whenever agents of types θ and ϕ meet, both become type $\theta + \phi$. This process is inductive over successive matching.

Setting for Information Percolation

Intuition: If the cross-sectional distribution of types is discrete, then the rate at which new agents of type θ are created is

$$2r \int \mu_t(\theta - z) \mu_t(z) dz = 2r(\mu_t * \mu_t)(\theta) \text{ a.s.}$$

This sort of application of the LLN for random matching is known as the Stosszahlansatz (Boltzmann), and has been shown rigorously only in discrete time (Duffie and Sun, *AAP*, 2007).

Solution for Cross-Sectional Distribution of Information

- ▶ The Boltzmann equation for the cross-sectional distribution μ_t of types is, for $\lambda = 2r$,

$$\frac{d}{dt} \mu_t = -\lambda \mu_t + \lambda \mu_t * \mu_t. \quad (1)$$

- ▶ **Standing assumption:** On the event $\{X = H\}$, the first moment of μ_0 is strictly positive, and μ_0 has a moment generating function $z \mapsto \int e^{z\theta} \mu_0(d\theta)$ that is finite on a neighborhood of $z = 0$.
- ▶ **Proposition (DGM, 2008).** The unique solution of (1) is the Wild sum

$$\mu_t = \sum_{n \geq 1} e^{-\lambda t} (1 - e^{-\lambda t})^{n-1} \mu_0^{*n}. \quad (2)$$

Sketch of Proof of Wild Sum

The ODE for the characteristic function $\varphi(\cdot, t)$ of μ_t ,

$$\frac{\partial \varphi(\mathbf{s}, t)}{\partial t} = -\lambda \varphi(\mathbf{s}, t) + \lambda \varphi^2(\mathbf{s}, t),$$

is solved by

$$\varphi(\mathbf{s}, t) = \frac{\varphi(\mathbf{s}, 0)}{e^{\lambda t}(1 - \varphi(\mathbf{s}, 0)) + \varphi(\mathbf{s}, 0)}.$$

This solution can be expanded as

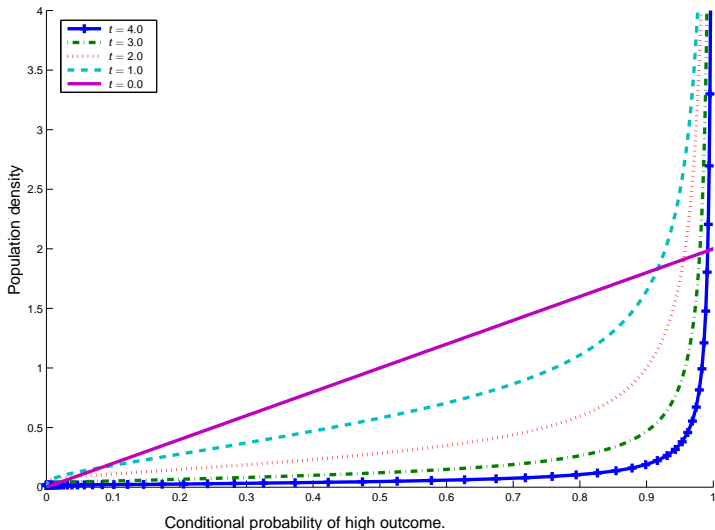
$$\varphi(\mathbf{s}, t) = \sum_{n \geq 1} e^{-\lambda t} (1 - e^{-\lambda t})^{n-1} \varphi^n(\mathbf{s}, t),$$

which is identical to the Fourier transform of the Wild sum (2).

Market Example

- ▶ Uninformed buyers of a contract promising X randomly select two informed sellers at intensity λ .
- ▶ A second-price auction allocates the trade to the lowest-bidding seller. (The Wallet Game.)
- ▶ In the unique symmetric equilibrium, sellers bid their posterior probabilities that X is high, revealing their types.

On the event $\{X = H\}$, the evolution of the cross-sectional population density of posterior probabilities of the event $\{X = H\}$.



Convergence Rate of Population Information

Let π_t be the cross-sectional distribution of posteriors at time t .

Definition: The rate of convergence of π_t to π_∞ is $\alpha > 0$ if there are constants κ_0 and κ_1 such that, for any b in $(0, 1)$,

$$e^{-\alpha t \kappa_0} \leq |\pi_t(0, b) - \pi_\infty(0, b)| \leq e^{-\alpha t \kappa_1}.$$

Proposition: π_t converges at rate λ to δ_0 on the event $\{X = L\}$ and to δ_1 on $\{X = H\}$.

Meetings of More than Two at a Time

- ▶ Groups of m agents are randomly matched. Because each agent is matched to others at rate r , the total annual quantity of attendance at meetings is $\lambda = mr$ a.s.
- ▶ The associated Boltzmann equation for the type distribution is

$$\frac{d}{dt}\mu_t = -\lambda\mu_t + \lambda\mu_t^{*m}.$$

- ▶ The solution is explicit as a Wild sum.

Wild Summation Solution

The unique solution of the Boltzmann equation for m -at-a-time matching is

$$\mu_t = \sum_{n \geq 1} a_{(m-1)(n-1)+1} e^{-\lambda t} (1 - e^{-(m-1)\lambda t})^{n-1} \mu_0^{*[(m-1)(n-1)+1]},$$

where $a_1 = 1$ and, for $n > 1$,

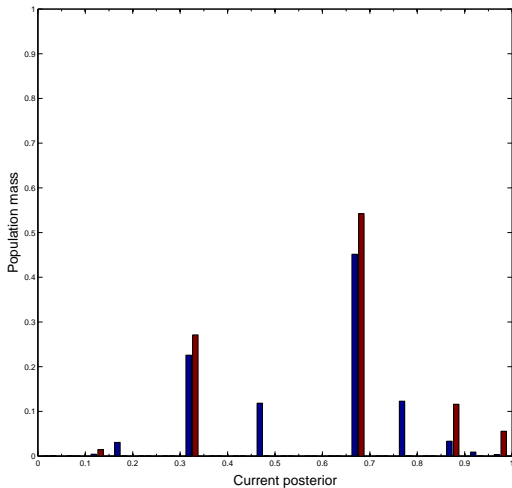
$$a_{(m-1)(n-1)+1} = \frac{1}{m-1} \left(1 - \sum_{\left\{ \begin{array}{l} i_1, \dots, i_{(m-1)} < n \\ \sum i_k = n+m-2 \end{array} \right\}} \prod_{k=1}^{m-1} a_{(m-1)(i_k-1)+1} \right).$$

Invariance of Convergence Rate to Group Size for a Given Total Rate of Meeting Attendance

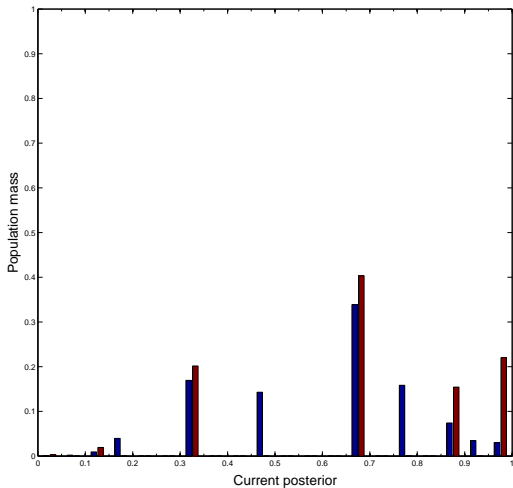
Proposition: For any group size m , the cross-sectional distribution π_t of posteriors converges at rate λ .

Malamud (2008) has extended this result to the case of groups of a random size.

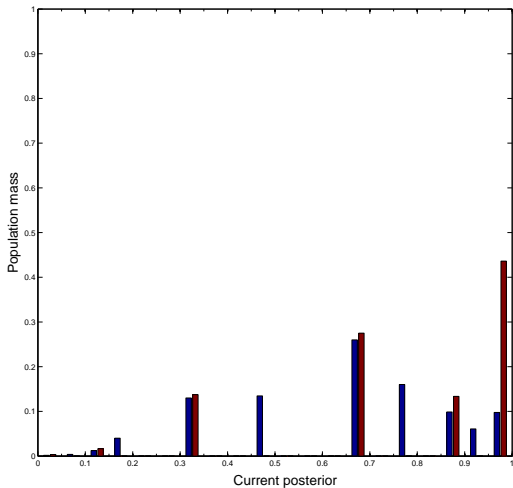
Groups of 2 (blue) versus Groups of 3 (red)



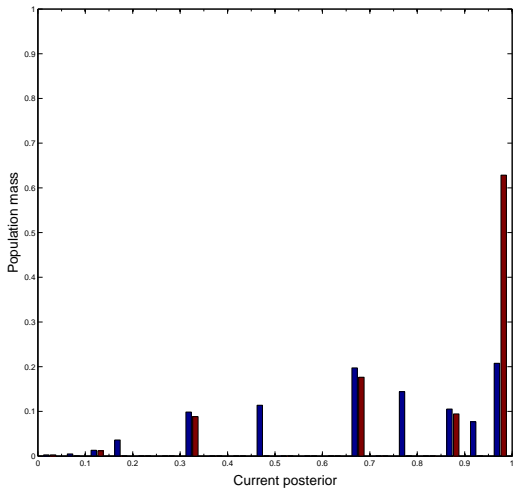
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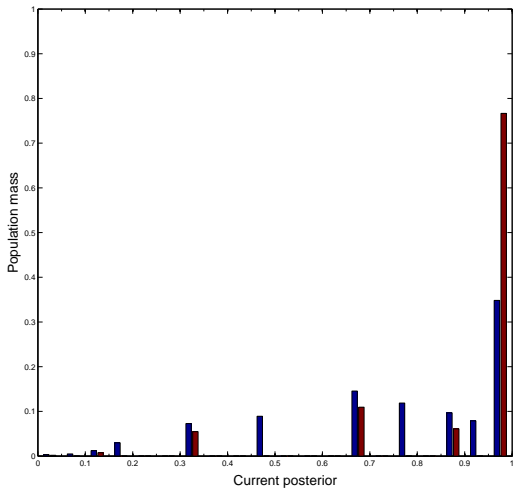
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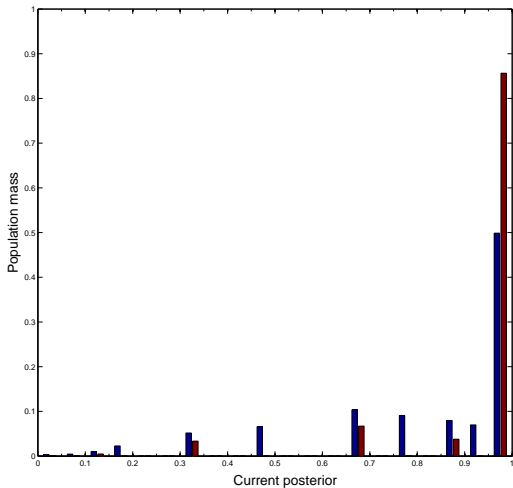
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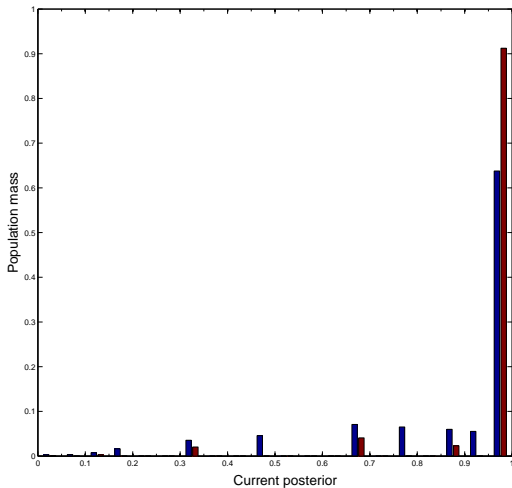
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Equilibrium Search Dynamics

With Manso and Malamud

- ▶ Signals and X are joint Gaussian, with $\text{corr}(X, s_i) = \rho$.
- ▶ Agents arrive at a rate proportional to the population size, and leave at exponentially distributed times, pairwise independently.
- ▶ Agents meet others at a mean rate proportional to the rate at which they choose to expend search costs.
- ▶ At entry, agent i receives N_{i0} signals, *iid* across agents.
- ▶ At exit, an agent chooses an action A , with cost $(X - A)^2$.
- ▶ The optimal exit action is $A = E(X | \mathcal{F}_{it})$, so the expected exit cost is the \mathcal{F}_{it} -conditional variance of X ,

$$\sigma_{it}^2 = v(N_{it}) \equiv \frac{1 - \rho^2}{1 + \rho^2(N_{it} - 1)}.$$

Information Transmission

- ▶ Agent i has current mean $E(X | \mathcal{F}_{it})$ and “precision” N_{it} .
- ▶ When agents i and j meet, their posterior precisions become $N_{it} + N_{jt}$. Their posterior means are given by the usual precision-weighted average of their priors.
- ▶ At constant meeting intensity λ , the cross-sectional precision distribution μ_t thus behaves as before, except for the effect of arrivals and departures.

Search Technology

- ▶ Random matching (Stosszahlansatz).
- ▶ Given current effort c by an agent, the mean rate of matching with someone from a unit mass of agents using search effort b is cb .
- ▶ The rate of cost of search effort c is $K(c)$, for $c \in [c_L, c_H]$.

Separability of Posterior Precision and Mean

Proposition. For any search-effort policy function $C : \mathbb{N} \rightarrow [c_L, c_H]$, the cross-sectional density f_t of precisions and posterior means of the agents is almost surely given by

$$f_t(n, y, \omega) = \mu_t(n) p_n(y | X(\omega)), \quad (3)$$

where μ_t is the unique solution of the Boltzmann equation for the evolution of the cross-sectional distribution of information precision and $p_n(\cdot | X)$ is the X -conditional Gaussian density of $E(X | s_1, \dots, s_n)$, for any n signals s_1, \dots, s_n .

Stationary Measure

In a stationary setting with search policy C , the precision distribution μ solves

$$0 = \eta(\gamma - \mu) + \mu^C * \mu^C - \mu^C \mu^C(\mathbb{N}),$$

where

- ▶ η is the mean replacement rate of agents.
- ▶ γ is the distribution of N_{i0} .
- ▶ μ^C is the effort-weighted measure, with $\mu^C(n) = C(n)\mu(n)$.

Stationary Measure

Lemma. Given any policy C , there is a unique measure μ satisfying the stationary-measure equation.

This measure μ is characterized in the paper, and under technical conditions is the pointwise limit of μ_t , invariant to μ_0 .

Optimality and Equilibrium

Given a policy C for other agents, agent i has the value function $V(\cdot)$ defined by

$$V(N_{it}) = \text{ess sup}_{\phi} E \left(-e^{-r(\tau-t)} v(N_{i\tau}) - \int_0^{\tau} e^{-r(u-t)} K(\phi_u) du \mid \mathcal{F}_{it} \right),$$

which is characterized by the associated HJB equation.

The policy C is an equilibrium if $\phi_t^* = C(N_{it})$ is optimal.

Trigger Policies

A trigger policy C^N , for some integer $N \geq 1$, is defined by

$$\begin{aligned} C_n^N &= c_H, & n < N, \\ &= c_L, & n \geq N. \end{aligned}$$

Information Sharing Opportunities

Proposition. Let μ^M and ν^N be the unique stationary measures corresponding to trigger policies C^M and C^N respectively. Let $\mu^{C,N}(n) = \mu^N(n)C^N(n)$ denote the associated search-effort-weighted measure. If $N > M$, then $\mu^{C,N}$ has first order stochastic dominance (FOSD) over $\mu^{C,M}$.

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This comparison result need not hold for non-trigger policies!

There exist cases with $B \leq C$ but μ^B having FOSD over μ^C .

Optimal Effort is Decreasing in Precision

Proposition. Suppose that K is increasing, convex, and differentiable. Then, given any population behavior (μ, C) , the optimal search effort policy function is decreasing in precision.

Proposition. Suppose that $K(c) = \kappa c$ for some scalar $\kappa > 0$. Then, given (μ, C) , some a trigger policy C^N that is optimal for all agents.

Existence of Equilibrium

Theorem. Suppose that $K(c) = \kappa c$ for some scalar $\kappa > 0$. Then there exists a trigger policy that is an equilibrium.

The Equilibrium Impact of a Search Subsidy

- ▶ A tax τ is charged to each agent entering the market
- ▶ The proceeds are used to subsidize search so that the search cost is $K_\delta(c) = (\kappa - \delta)c$.

Proposition. If C^N is an equilibrium with subsidy δ , then for any $\delta' \geq \delta$, there exists some $N' \geq N$ such that $C^{N'}$ is an equilibrium with subsidy δ' .

Example

1. For some integer $N > 1$, $\gamma(0) = 1/2$, $\gamma(N) = 1/2$, and $c_L = 0$.
2. It is possible to choose parameters so that, given market conditions (μ^N, C^N) , agents slightly prefer policy C^0 over C^N .
3. We can choose the subsidy rate δ so that, given market conditions (μ^N, C^N) , agents strictly prefer C^N to C^0 .
4. For sufficiently large N all agents have strictly higher indirect utility.

The Equilibrium Impact of Public Information

Agents are given $M \geq 1$ additional public signals at entry.

Proposition If C^N is an equilibrium with M public signals, then for any $M' \leq M$, there exists some $N' \geq N$ such that $C^{N'}$ is an equilibrium with M' public signals.

We provide examples with strict dominance.

Example

1. Suppose, for some integer $N > 1$, that $\pi_0 = 1/2$, $\pi_N = 1/2$, and $c_L = 0$.
2. Choose parameters so that, given market conditions (μ^N, C^N) agents are indifferent between policies C^N and C^0 .
3. Give each agent $M = 1$ public signal at entry.
4. All agents strictly prefer C^0 to C^N
5. For sufficiently large N all agents have strictly lower indirect utility.

Conclusion

- ▶ Model of social learning with **endogenous search intensity**.
- ▶ Social learning may slow down or even collapse:
 - coordination problems.
 - externality problems.
- ▶ Two policy interventions:
 - search subsidy.
 - education at entry.

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